

Segmentation

- (1) How are relocation and protection implemented in Pentium architecture? Consider the roles of both segmentation and paging.
- (2) How is segmentation different from paging? Why was each technique invented?
- (3) Using either Multics or Pentium architecture as an example, explain how segmentation is used in enforcing protection?
- (4) How is a virtual address converted to a physical address, considering both segmentation and paging in (a) Multics and (b) Pentium architectures?
- (5) In paged-segmentation implementations: Multics has one page table per segment, whereas Pentium has one page table for multiple segments. Why the difference? Which one is better? Why?
- (6) OR
- (7) How many page tables are there per segment in Multics and Pentium. (b) Which one (Multics/Pentium) is better? Why?
- (8) What problem does segmentation solve that paging doesn't solve? What problem does paging solve that segmentation doesn't solve?
- (9) Which of the following memory designs can cause internal fragmentation only, external fragmentation only, both, or neither? Briefly explain why.
 - A. Pure paging
 - B. Pure segmentation
 - C. Paging with Segmentation
 - D. Using superpages of the same size (no base pages or any other superpage size)
 - E. Using a mix of superpages of different sizes
- (10) Considering memory protection, explain how the operating system ensures that user-level processes don't access kernel-level memory?
- (11) Consider segment-level protection and CPU privilege levels in x86. Explain how the MMU, together with OS, determines whether a code executing on the CPU, has the permissions to access a given memory address?